

Linear Probing: The Probable Largest Search Time Grows Logarithmically with the Number of Records

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The worst-case performance of a linear probing algorithm is studied under assumption that each of m locations can contain at most one record. It is shown that the length of the longest sequence of probes needed to insert (locate) a record grows, in probability, as $c^{-1} \log m$ if m and n , the total number of records, approach ∞ in such a way that the load factor $a = n/m$ is bounded away from 0 and 1, here $c = a - 1 - \log a$. The argument is based on the study of the size distribution of the longest block of occupied locations. © 1987 Academic Press, Inc.

1. INTRODUCTION: RESULTS

Linear probing is a hashing technique in which a record x hashed to an already occupied cell $h(x)$ is sent to the next empty cell modulo m . ($h(\cdot)$ is a hashing function, and m is the size of the table.) This empty location is found via the sequential search beginning at the cell $h(x)$, and continuing at the cell 1 if no cell with index $\geq h(x)$ is empty. (It is convenient to think that the cells $1, \dots, m$ form an oriented cycle.) As usual, we assume that the sequence of records x_1, \dots, x_n ($n \leq m$) and the hash function $h(\cdot)$ are such that $h(x_1), \dots, h(x_m)$ are independent, uniformly distributed over $\{1, \dots, m\}$. Denote T_j the random number of probes needed to insert (locate) the record x_j , $1 \leq j \leq n$. The distribution and the moments of T_j were studied in [4, 5, 6]. (A more general scheme with each location capable to accommodate more than one record was analyzed in [1].)

Our goal is to study $L_n = \max_{1 \leq j \leq n} T_j$, which is interpreted as the worst-case search time for the first n records. To this end, we introduce another random variable B_n , the length of the longest block of cells

occupied by the records x_1, \dots, x_n . It should be clear that

$$L_n \leq B_{n-1} + 1. \quad (1.1)$$

Suppose that $m, n \rightarrow \infty$ and that the load factor $a (= n/m)$ remains bounded away from 0 and 1. Introduce a parameter

$$c = a - 1 - \log a,$$

it is clear that $c \in (0, \infty)$ and is bounded away from 0 and ∞ . We shall prove

THEOREM. *Define a random variable B_n^* by the relation*

$$B_n = c^{-1} \left[\log m - \frac{3}{2} \log \log m + B_n^* \right] - 1.$$

Then, B_n^ is double exponentially distributed in the limit. More precisely, if $y = O(1)$ is such that*

$$b = c^{-1} \left[\log m - \frac{3}{2} \log \log m + y \right] - 1$$

is a positive integer then

$$P(B_n^* < y) - \exp(-\lambda e^{-y}) \rightarrow 0.$$

Here

$$\lambda = c^{3/2} (a^{-1} - 1) \left[(1 - e^{-c})(2\pi)^{1/2} \right]^{-1}. \quad (1.2)$$

Thus, in probability,

$$B_n = c^{-1} \left[\log m - \frac{3}{2} \log \log m \right] + O(1). \quad (1.3)$$

COROLLARY. *For every $\delta \in (\frac{1}{2}, 1)$,*

$$c^{-1} \log m + O_p(\log^\delta m) \leq L_n \leq c^{-1} \left[\log m - \frac{3}{2} \log \log m \right] + O_p(1). \quad (1.4)$$

Here $O_p(\log^\delta m), O_p(1)$ stand for random terms Δ_1, Δ_2 such that $\Delta_1/\omega(n)\log^\delta m, \Delta_2/\omega(n) \rightarrow 0$ in probability, for every $\omega(n) \rightarrow \infty$.

Proof. According to (1.1), we only have to prove the lower estimate. Introduce $n_1 = n - [n \log^{\delta-1} m]$, then

$$a_1 = n_1/m = a + O(\log^{\delta-1} m), \quad c_1 = c + O(\log^{\delta-1} m),$$

and

$$c_1^{-1}[\log m - \frac{3}{2}\log \log m] = c^{-1}\log m + O(\log^\delta m).$$

By (1.3),

$$P(B_{n_1} < c^{-1}\log m - \gamma \log^\delta m) \rightarrow 0, \tag{1.5}$$

if γ is sufficiently large. If, on the other hand, the hash sequence $h(x_1), \dots, h(x_{n_1})$ is such that $B_{n_1} \geq c^{-1}\log m - \gamma \log^\delta m$, then the conditional probability that $L_n < c^{-1}\log m - 2\gamma \log^\delta m$ is at most

$$(1 - [\gamma \log^\delta m]/m)^{n-n_1} \leq \exp(-2^{-1}\gamma a \log^{2\delta-1} m) = o(1). \tag{1.6}$$

(Indeed, if at least one of the records x_{n_1+1}, \dots, x_n is hashed into one of $[\gamma \log^\delta m]$ first cells of a maximal block corresponding to the records x_1, \dots, x_{n_1} then the number of probes for this record is at least $c^{-1}\log m - 2\gamma \log^\delta m$.) The relations (1.5), (1.6) together imply that, with probability approaching 1,

$$L_n \geq c^{-1}\log m - 2\gamma \log^\delta m.$$

How sharp are the estimates (1.4)? Using the explicit formulae for $P(T_j > k)$, [4, 5], we shall also prove that

$$L_n \leq c^{-1}[\log m - (5/2)\log \log m] + O_p(1). \tag{1.7}$$

We conjecture that in this relation \leq can be replaced by $=$.

The reader interested in the largest search time analysis of other hashing methods is referred to [3] (uniform probing, random probing, direct chaining), [2] (direct chaining for a nonuniform distribution of records), and [8, 9] (coalesced hashing).

2. PROOFS

We begin with

LEMMA 1. For every $b \geq 1$,

$$P(B_n < b) = [n!/m^{n-1}(m-n)] \text{coeff}_{x^m} [t_b(x)]^{m-n}, \tag{2.1}$$

where

$$t_b(x) = \sum_{k=1}^b x^k k^{k-1}/k!.$$

Notes. (1) $t_b(x)$ is the b th partial sum of the exponential generating function of rooted trees. (2) Setting $m = \alpha$, $m - n = \beta$, and $b = \infty$ we get from (2.1) that

$$\text{coeff}_{x^\alpha} [t(x)]^\beta = \beta \alpha^{\alpha - \beta - 1} / (\alpha - \beta)!,$$

a relation crucial for enumeration of (labelled) trees and forests, [7].

Proof of the Lemma. Denote $S = \{1, \dots, m\}$. Let $A(s^n)$ be the set of n occupied cells if the algorithm is applied to s^n , $s^n = (x_1, \dots, x_n) \in S^n$. Since

$$P(h(x_1) = s_1, \dots, h(x_n) = s_n) = m^{-n}, \quad \forall s^n \in S^n,$$

the problem is to evaluate the number of s^n such that $A(s^n)$ does not contain a block of size $\geq b$.

Fix $\nu < m$ and $i \leq m$. According to [4, 5], the number of s^ν such that

$$A(s^\nu) = \{i, i + 1, \dots, i + \nu - 1\}$$

(addition modulo m) equals

$$f(\nu) = (\nu + 1)^{\nu - 1}. \quad (2.2)$$

Now, given nonnegative integers k_1, \dots, k_n satisfying

$$\sum_{j=1}^n j k_j = n, \quad k_j = 0 \quad \text{for } j \geq b, \quad (2.3)$$

denote $H_{mn}(k)$ the number of all s^n such that $A(s^n)$ contains k_j blocks of size j , $i \leq j \leq n$. It turns out that

$$H_{mn}(k) = mn!(m - n - 1)_{(|k|-1)} \prod_{j=1}^n [f(j)/j!]^{k_j} / k_j!, \quad (2.4)$$

where

$$|k| = \text{def} \sum_{j=1}^n k_j, \quad (\alpha)_\beta = \text{def} \alpha(\alpha - 1) \cdots (\alpha - \beta + 1).$$

To prove this formula, we describe a procedure which generates all such sequences s^n . First, we partition the set $\{1, \dots, n\}$ into k_1 subsets of size 1,

k_2 subsets of size 2, ..., and k_n subsets of size n , which can be done in

$$N_1 = n! \left/ \prod_{j=1}^n (j!)^{k_j} k_j! \right.$$

ways. Given a partition \mathcal{P} , each of its $|k|$ subsets is assigned to a block of cells from the oriented cycle of m cells, so that the cardinality of a subset equals the size of the corresponding block, and all the blocks are disjoint. To define such an assignment we (1) select one of $N_{21} = (|k| - 1)!$ possible cyclic orders of the blocks, (2) select one of $N_{22} = m$ possible positions for the block corresponding to the subset which contains, say, the element 1, and finally (3) select the sizes of all $|k|$ gaps between the neighboring blocks. Since the total length of the gaps is $m - n$, the selection (3) can be done in the number of ways N_{23} equal to the number of integer solutions of

$$\sum_{t=1}^{|k|} \Delta_t = m - n, \quad \Delta_t \geq 1, \quad 1 \leq t \leq |k|,$$

whence

$$N_{23} = \binom{m - n - 1}{|k| - 1}.$$

After this, it remains to select for every subset $\{\alpha_1, \dots, \alpha_\nu\}$ of the partition \mathcal{P} a sequence $s^\nu \in S^\nu$ such that $A(s^\nu)$ coincides with the block of cells that subset $\{\alpha_1, \dots, \alpha_\nu\}$ is assigned to. Since it can be done in $f(\nu)$ ways, the total selection can be done in

$$N_3 = \prod_{j=1}^n [f(j)]^{k_j}$$

ways. Clearly,

$$H_{mn}(k) = N_1 \cdot (N_{21}N_{22}N_{23}) \cdot N_3,$$

which is equivalent to (2.4).

Denote D_n the random number of blocks of all sizes. By (2.4),

$$P(D_n = d, B_n < b) = m^{-n} [mn!(m - n - 1)_{(d-1)}] P(n, d, b), \quad (2.5)$$

where

$$P(n, d, b) = \sum_k \prod_{j=1}^n [f(j)/j!]^{k_j} / k_j!,$$

k satisfies (2.3) and such that $|k| = d$. But

$$\begin{aligned}
 & 1 + \sum_{n \geq 1, d \geq 1} P(n, d, b) x^n y^d \\
 &= 1 + \sum_{n \geq 1, d \geq 1} \sum_k \prod_{j=1}^n [y x^j f(j)/j!]^{k_j} / k_j! \\
 &= \sum_{k_1, \dots, k_{b-1} \geq 0} \prod_{j=1}^{b-1} [y x^j f(j)/j!]^{k_j} / k_j! \\
 &= \prod_{j=1}^{b-1} \left\{ \sum_{k_j \geq 0} [y x^j f(j)/j!]^{k_j} / k_j! \right\} \\
 &= \exp \left[y \sum_{j=1}^{b-1} x^j f(j)/j! \right].
 \end{aligned}$$

Hence

$$\begin{aligned}
 P(n, d, b) &= \text{coeff}_{x^n} \text{coeff}_{y^d} \exp \left[y \sum_{j=1}^{b-1} x^j f(j)/j! \right] \\
 &= \text{coeff}_{x^n} \left[\sum_{j=1}^{b-1} x^j f(j)/j! \right]^d / d!.
 \end{aligned} \tag{2.6}$$

So, (see (2.5)),

$$\begin{aligned}
 P(B_n < b) &= [n! / m^{n-1} (m-n)] \\
 &\quad \times \text{coeff}_{x^n} \left[\sum_{d=0}^{m-n} \binom{m-n}{d} \left(\sum_{j=1}^{b-1} x^j f(j)/j! \right)^d \right],
 \end{aligned}$$

where, according to (2.2),

$$\begin{aligned}
 \sum_{d=0}^{m-n} \binom{m-n}{d} \left(\sum_{j=1}^{b-1} x^j f(j)/j! \right)^d &= \left[1 + \sum_{j=1}^{b-1} x^j (j+1)^{j-1} / j! \right]^{m-n} \\
 &= x^{n-m} \left(\sum_{j=1}^b x^j j^{j-1} / j! \right)^{m-n}.
 \end{aligned}$$

The last two relations are equivalent to (2.1).

This lemma and contour integration enables us to prove

LEMMA 2. Let $n, m \rightarrow \infty$. Suppose that an integer $b \rightarrow \infty$ in such a way that

$$mb^{-5/2}e^{-cb} = o(1), \tag{2.7}$$

(recall that $c = a - 1 - \log a$, $a = n/m$). Then

$$P(B_n < b) \sim \exp(-\delta m \beta^{-3/2} e^{-c\beta}), \quad \beta = b + 1,$$

where

$$\delta = (a^{-1} - 1) \left[(1 - e^{-c})(2\pi)^{1/2} \right]^{-1}. \tag{2.8}$$

Applying Lemma 2 to

$$b = c^{-1} [\log m - (3/2) \log \log m + y] - 1, \quad y = O(1),$$

we immediately obtain the proof of theorem.

Proof of Lemma 2. From (2.1) and Cauchy's formula, it follows that

$$P(B_n < b) = [n!/m^{n-1}(m-n)](2\pi i)^{-1} I, \tag{2.9}$$

$$I = \int_C z^{-(m+1)} [t_b(z)]^{m-n} dz, \tag{2.10}$$

where C is a circle with center at $z = 0$. To estimate I , we use the saddle-point method. To this end, we select the radius of C equal to r , which is a minimum point of the function $F(x) = x^{-m} [t_b(x)]^{m-n}$, $x \in (0, \infty)$. (Since n/m is bounded away from 0 and 1, the minimum is achieved if b is sufficiently large.) So, r is a root of an equation

$$xt'_b(x)/t_b(x) = (1 - a)^{-1}, \quad a = n/m. \tag{2.11}$$

This equation has precisely one root, since by the Cauchy-Schwartz inequality

$$\begin{aligned} \frac{d}{dx} [xt'_b(x)/t_b(x)] &= [xt_b^2(x)]^{-1} \left[\left(\sum_{k=1}^b x^k k^2 c_k \right) \right. \\ &\quad \left. \times \left(\sum_{k=1}^b x^k c_k \right) - \left(\sum_{k=1}^b x^k k c_k \right)^2 \right] > 0, \\ &(c_k = k^{k-1}/k!). \end{aligned}$$

It is possible to estimate r rather sharply for large bs . Namely, it is known

[7] that the series $\sum_{j=1}^{\infty} x^j j^{-1}/j!$ converges for $|x| \leq e^{-1}$, and its sum $t(x)$ is the solution of an equation $t = xe^t$, so that, in particular,

$$xt'(x)/t(x) = [1 - t(x)]^{-1}, \quad |x| \leq e^{-1}. \quad (2.12)$$

In view of these relations, $\bar{r} = ae^{-a}$ is the solution of an equation

$$xt'(x)/t(x) = (1 - a)^{-1}, \quad (2.13)$$

and $t(\bar{r}) = a$. Since $t_b(x)$, $t'_b(x)$ converge (resp. to $t(x)$, $t'(x)$) uniformly in every interval $|x| \leq \xi$, $\xi < e^{-1}$, and

$$[xt'(x)/t(x)]'|_{x=\bar{r}} = e^a(1 - a)^{-3} > 0,$$

we have then

$$r = \bar{r} + \Delta, \quad \Delta = o(1), \quad (b \rightarrow \infty). \quad (2.14)$$

In the Appendix, we show how this estimate leads, via the estimations of $t(r) - t_b(r)$, $t'(r) - t'_b(r)$, to

$$\Delta = \theta_1 \beta^{-1/2} \rho^\beta + \theta_2 \beta^{-3/2} \rho^\beta + O(\beta^{-5/2} \rho^\beta), \quad (2.15)$$

$$t_b(r) = a + e^a(1 - a)^{-1} \Delta + \theta_3 \beta^{-3/2} \rho^\beta + O(\beta^{-5/2} \rho^\beta), \quad (2.16)$$

here

$$\theta_1 = (1 - a)^3 [ae^a(2\pi)^{1/2}(1 - \rho)]^{-1}, \quad (2.17)$$

$$\theta_2 = -\theta_1 [12^{-1}(1 + 5\rho)(1 - \rho)^{-1} + (1 - a)^{-1}], \quad (2.18)$$

$$\theta_3 = -(2\pi)^{-1/2}(1 - \rho)^{-1}, \quad (2.19)$$

$\beta = b + 1$ and $\rho = er$. (Observe that $e\bar{r} = ae^{1-a} < 1$.) According to (2.14)–(2.19), we have then

$$\begin{aligned} F(r) &= r^{-m} [t_b(r)]^{m-n} = (a^{m-n}/\bar{r}^m) \exp \left\{ -m \log(1 + \Delta/r) \right. \\ &\quad \left. + (m - n) \log \left[1 + \Delta(1 - a)^{-1}/\bar{r} + a^{-1} \theta_3 \beta^{-3/2} \rho^\beta + O(\beta^{-5/2} \rho^\beta) \right] \right\} \\ &= (a^{m-n}/\bar{r}^m) \exp \left[a^{-1} \theta_3 (m - n) \beta^{-3/2} \rho^\beta + O(m \beta^{-5/2} \rho^\beta) \right] \\ &\sim (a^{m-n}/\bar{r}^m) \exp(-\delta m \beta^{-3/2} e^{-c\beta}), \quad c = -\log(e\bar{r}); \end{aligned} \quad (2.20)$$

(by (2.15), $\rho^\beta = (e\bar{r})^\beta + O(\beta^{1/2}(e\bar{r})^{2\beta})$).

Now, set in (2.10) $z = re^{i\phi}$, $\phi \in (-\pi, \pi]$ and write

$$\begin{aligned} I &= \int_C F(z) z^{-1} dz = i \int_{-\pi}^{\pi} F(re^{i\phi}) d\phi \\ &= i \int_{|\phi| \leq \phi_0} + i \int_{|\phi| > \phi_0} = iI_1 + iI_2, \end{aligned}$$

where $\phi_0 = m^{-5/12}$. Consider I_1 . Since $F'(r) = 0$, we have

$$\begin{aligned} F(re^{i\phi}) &= \exp[\log F(re^{i\phi})] \\ &= F(r) \exp[-\alpha(r)\phi^2/2 + O(m\phi_0^3)], \end{aligned}$$

where (see (2.10), (2.11)),

$$\begin{aligned} \alpha(r) &= r^2 [\log F(z)]''|_{z=r} = m + (m-n) [r^2 t_b''(r)/t_b(r) - (1-a)^{-2}] \\ &\sim m [1 + (1-a)\bar{r}^2 t''(\bar{r})/t(\bar{r}) - (1-a)^{-1}] = ma(1-a)^{-2}. \end{aligned}$$

($t''(r)$ is obtained by differentiating both sides of (2.12) at $x = \bar{r}$.) Since $m\phi_0^3 \rightarrow 0$, $m\phi_0^2 \rightarrow \infty$, we conclude that

$$\begin{aligned} I_1 &\sim F(r) \int_{|\phi| \leq \phi_0} \exp[-\alpha(r)\phi^2/2] d\phi \\ &\sim F(r) [2\pi(1-a)^2 m^{-1} a^{-1}]^{1/2}. \end{aligned} \quad (2.21)$$

It remains to estimate I_2 . We shall prove in the Appendix that, for all ϕ ,

$$|t_b(re^{i\phi})| \leq a \exp[r(\cos \phi - 1)/2 + O(\beta^{-3/2} e^{-c\beta})]. \quad (2.22)$$

Invoking (2.7) and an inequality $\cos \phi - 1 \leq -\omega\phi^2$, ($\phi \in (-\pi, \pi)$), $\omega > 0$, we obtain then for $|\phi| > \phi_0$, and large enough m ,

$$\begin{aligned} |t_b(re^{i\phi})| &\leq a \exp[-ae^{-a}\omega\phi^2 + o(m^{-1}\log m)] \\ &\leq a \exp[-ae^{-a}\omega m^{-5/6} + o(m^{-5/6})] \\ &\leq a \exp(-\omega_0 m^{-5/6}), \quad \omega_0 > 0. \end{aligned} \quad (2.23)$$

Using (2.23), and also (2.15), we arrive at

$$\begin{aligned} |I_2| &\leq 2\pi(a^{m-n}/r^m)\exp[-\omega_0(1-a)m^{1/6}] \\ &= (a^{m-n}/\bar{r}^m)\exp[-\omega_0(1-a)m^{1/6} + O(m\beta^{-1/2}e^{-c\beta})], \end{aligned}$$

which, together with (2.7), (2.20), and (2.21), implies that

$$|I_2| = o(I_1). \quad (2.24)$$

Combination of (2.9), (2.10), (2.20), (2.21), (2.24) yields, after numerous cancellations,

$$\begin{aligned} P(B_n < b) &\sim \left[(2\pi n)^{1/2} (n/e)^n / (m^{n-1}(m-n)) \right] (2\pi)^{-1} \\ &\quad \times (a^{m-n}/\bar{r}^m) \left[2\pi(1-a)^2 m^{-1} a^{-1} \right]^{1/2} \\ &\quad \times \exp(-\delta m \beta^{-3/2} e^{-c\beta}) \\ &= \exp(-\delta m \beta^{-3/2} e^{-c\beta}). \quad (!) \end{aligned}$$

The lemma is proven.

Let us finally prove that

$$L_n \leq c^{-1} \left[\log m - \frac{5}{2} \log \log m \right] + O_p(1). \quad (2.25)$$

Suppose that an integer b is such that

$$b = c^{-1} \left[\log m - \frac{5}{2} \log \log m + \omega(m) \right], \quad (2.26)$$

where $\omega(m) \rightarrow \infty$. Denote $n_1 = \lfloor n/2 \rfloor$. By the theorem,

$$\begin{aligned} P(L_n > b) &= o(1) + P\left(\max_{n_1 \leq \nu \leq n} T_\nu > b \right) \\ &\leq o(1) + \sum_{\nu=n_1}^n P(T_\nu > b) = o(1) + \sum_{\nu=n_1}^n \sum_{k=b+1}^{\nu} P(T_\nu = k). \end{aligned} \quad (2.27)$$

Here [4, 5],

$$\begin{aligned} P(T_\nu = k) &= m^{-\nu+1} \sum_{j=k}^{\nu} \binom{\nu-1}{j-1} j^{j-2} (m-j)^{\nu-1-j} (m-\nu) \\ &=_{\text{def}} \sum_{j=k}^{\nu} P(\nu, j). \end{aligned} \quad (2.28)$$

Furthermore, for $k \leq j < \nu$,

$$\begin{aligned} P(\nu, j + 1)/P(\nu, j) &= \left(\frac{\nu - j}{m - j - 1} \right) (1 + j^{-1})^{j-1} (1 - (m - j)^{-1})^{\nu-1-j} \\ &\leq \left(\frac{\nu - j}{m - j - 1} \right) \exp\left(1 - \frac{\nu - 1 - j}{m - j}\right) \\ &= [\eta + O(m^{-1})] e^{1-\eta}, \end{aligned} \tag{2.29}$$

where

$$\eta = \frac{\nu - 1 - j}{m - j} \leq n/m = a.$$

Since $ae^{1-a} < 1$ and bounded away from 1, (2.28) implies that

$$P(T_\nu = k) = O(P(\nu, k)),$$

and, (cf. (2.28) and the innermost sum in (2.27)),

$$P(L_n > b) = O\left(\sum_{\nu=n_1}^n P(\nu, b + 1)\right) + o(1). \tag{2.30}$$

Here, by Stirling's formula,

$$\begin{aligned} P(\nu, b + 1) &\leq \text{const. } m^{-b} \binom{\nu - 1}{b} b^{b-1} \exp(-b\nu/m) \\ &\leq \text{const. } b^{-3/2} m^{-b} \exp(-b\nu/m) \frac{(\nu - 1)^{\nu-1}}{(\nu - 1 - b)^{\nu-1-b}} \\ &\leq \text{const. } b^{-3/2} m^{-b} \exp[-b\nu/m + b(1 + \log \nu)] \\ &= \text{const. } b^{-3/2} \exp[-bc(\nu/m)], \end{aligned}$$

where

$$c(x) = x - 1 - \log x, \quad x \in (0, 1].$$

(Note that $c(a) = c$.) Since $c'(x) < 0$, $c''(x) > 0$ for $x \in (0, 1)$,

$$\begin{aligned} \sum_{\nu=n}^n P(\nu, b + 1) &\leq \text{const. } b^{-3/2} e^{-cb} \sum_{\nu=n_1}^n \exp[m^{-1}bc'(a)(n - \nu)] \\ &\leq \text{const. } m^{-1} b^{-5/2} e^{-cb}. \end{aligned}$$

Hence (see (2.26), (2.30)),

$$P(L_n > b) = O(\exp(-\omega(m))) = o(1),$$

and the relation (2.25) is proven.

APPENDIX

Show first that uniformly in any disk $|z| \leq R$, $R < e^{-1}$,

$$\sum_{k \geq \beta} k^{k-1} z^k / k! = (2\pi)^{-1/2} \beta^{-3/2} \xi^\beta (1 - \xi)^{-1} + O(\beta^{-5/2} |\xi|^\beta), \quad (1)$$

$$\begin{aligned} \sum_{k \geq \beta} k^k z^k / k! &= (2\pi)^{-1/2} \beta^{-1/2} \xi^\beta (1 - \xi)^{-1} \\ &\quad - 12^{-1} (2\pi)^{-1/2} \beta^{-3/2} \xi^\beta (1 + 5\xi) (1 - \xi)^{-2} \\ &\quad + O(\beta^{-5/2} |\xi|^\beta), \end{aligned} \quad (2)$$

$\xi = ze$. Really, using the Stirling series for $\log k!$, we have

$$k^{k-1} z^k / k! = (2\pi)^{-1/2} k^{-3/2} \xi^k + O(k^{-5/2} |\xi|^k), \quad (3)$$

$$k^k z^k / k! = (2\pi)^{-1/2} k^{-3/2} \xi^k - 12^{-1} (2\pi)^{-1/2} k^{-3/2} \xi^k + O(k^{-5/2} |\xi|^k). \quad (4)$$

Together with

$$k^{-1/2} = \beta^{-1/2} - 2^{-1} \beta^{-3/2} (k - \beta) + O(\beta^{-5/2} (k - \beta)^2), \quad k \geq \beta,$$

$$k^{-3/2} = \beta^{-3/2} + O(\beta^{-5/2} (k - \beta)), \quad k \geq \beta,$$

the relations (3)–(4) imply (1)–(2).

Proof of (2.15)–(2.16). According to (1)–(2),

$$t_b(r) = t(r) + \omega_1 \beta^{-3/2} \rho^\beta + O(\beta^{-5/2} \rho^\beta), \quad (5)$$

$$rt'_b(r) = rt'(r) + \omega_1 \beta^{-1/2} \rho^\beta + \omega_2 \beta^{-3/2} \rho^\beta + O(\beta^{-5/2} \rho^\beta), \quad (6)$$

where

$$\omega_1 = -(2\pi)^{-1/2} (1 - \rho)^{-1}, \quad \omega_2 = 12^{-1} (2\pi)^{-1/2} (1 + 5\rho) (1 - \rho)^{-2},$$

$\beta = b + 1$, and $\rho = er$. Further, since $r = \bar{r} + \Delta$, $\Delta = o(1)$, $\bar{r} = ae^{-a}$ and

$t(\bar{r}) = a$, we have (see (2.12)–(2.13)),

$$t(r) = t(\bar{r}) + t'(\bar{r})\Delta + O(\Delta^2) = a + e^a(1 - a)^{-1}\Delta + O(\Delta^2), \tag{7}$$

$$rt'(r) = t(r)[1 - t(r)]^{-1} = a(1 - a)^{-1} + e^a(1 - a)^{-3}\Delta + O(\Delta^2). \tag{8}$$

Therefore (see (2.11), (5–8)),

$$\begin{aligned} & a(1 - a)^{-1} + e^a(1 - a)^{-3}\Delta + O(\Delta^2) + \omega_1\beta^{-1/2}\rho^\beta + \omega_2\beta^{-3/2}\rho^\beta \\ &= a(1 - a)^{-1} + e^a(1 - a)^{-2}\Delta + \omega_1(1 - a)^{-1}\beta^{-3/2}\rho^\beta \\ &+ O(\beta^{-5/2}\rho^\beta), \end{aligned}$$

which yields (see (2.17)–(2.19)),

$$\Delta = \theta_1\beta^{-1/2}\rho^\beta + \theta_2\beta^{-3/2}\rho^\beta + O(\beta^{-5/2}\rho^\beta).$$

Also, the relation (2.16) follows immediately from (5), (7), so that $\theta_3 = \omega_1$.

Proof of (2.22). Since by (1)

$$|t_b(re^{i\phi})| \leq |t(re^{i\phi})| + O(\beta^{-3/2}e^{-c\beta}),$$

it suffices to show that, for $|z| \leq e^{-1}$,

$$|t(z)| \leq t(|z|)\exp[(\operatorname{Re} z - |z|)/2],$$

(cf. [10], an inequality (2.16)). To this end, fix a number $c > 0$, and write

$$t(z) = e^{cz} [e^{-cz}t(z)] = e^{cz} \sum_{n=1}^{\infty} a_n z^n, \tag{9}$$

where

$$a_n = \sum_{j=0}^{n-1} a_{nj}, \quad a_{nj} = (-1)^j c^j (n - j)^{n-j-1} / j!(n - j)!.$$

Further, for $1 \leq j \leq n - 2$,

$$\begin{aligned} |a_{n,j+1}| / |a_{nj}| &= c(j + 1)^{-1} [1 - (n - j)^{-1}]^{n-j-2} \\ &\leq ce^{-1}(j + 1)^{-1} [1 - (n - j)^{-1}]^{-2} \leq 2ce^{-1}. \end{aligned}$$

If $\chi = 2ce^{-1} < 1$ then

$$\sum_{j=1}^{n-1} |a_{nj}| \leq a_{n0}\chi(1 - \chi)^{-1}$$

and

$$a_n \geq a_{n_0} - \sum_{j=1}^{n-1} |a_{nj}| \geq a_{n_0} [1 - \chi(1 - \chi)^{-1}].$$

Since $a_{n_0} = n^{n-1}/n! > 0$, all the coefficients are therefore positive provided that $\chi < \frac{1}{2}$, or that $c < e/4$. Take $c = \frac{1}{2}$; as $a_n > 0$, $n = 1, 2, \dots$, it follows from (9) that

$$\begin{aligned} |t(z)| &= |e^{z/2}| \cdot \left| \sum_{n=1}^{\infty} a_n z^n \right| \leq \exp(\operatorname{Re} z/2) \cdot \sum_{n=1}^{\infty} a_n |z|^n \\ &= \exp[(\operatorname{Re} z - |z|)/2] e^{|z|/2} \cdot \sum_{n=1}^{\infty} a_n |z|^n \\ &= \exp[(\operatorname{Re} z - |z|)/2] t(|z|). \end{aligned}$$

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