

# Divide and Conquer

Lecture outline

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Call the roll. Preliminary questions?

## 1 Brute force

- Exhaustive exploration
- Examples: linear search, sort
- Often possible, always consider before discarding the option
- Sometimes the only solution
- Sometimes very small constant, ok for small input sizes

## 2 Divide and Conquer (4.1, 4.2)

- Historical origin (*Divide et impera*, attributed to Philip II of Macedon, used in politics, in military strategy)
- In computer science: divide, process recursively, merge
- Often allows (but not always) saving time

### 2.1 Example: binary search

### 2.2 Example: matrix multiplication

- Standard algorithm
- Naïve divide and conquer, superficial analysis, recurrence formula
- Mention Strassen algorithm and its recurrence formula

$$T(n) = \begin{cases} \Theta(1) & \text{if } n = 1; \\ 7T(\frac{n}{2}) + \Theta(n^2) & \text{if } n > 1. \end{cases}$$

### 3 Analyzing Recursive Algorithms (4.4)

- Simplifying assumption of powers of 2 (or of another number), not critical
- Recurrence trees for binary search, naïve matrix multiplication, Strassen ( $\log_2 7 \approx 2.807$ )
- Practical use of Strassen algorithm
- Mention of modern matrix multiplications (Coppersmith–Winograd and improvements, 2.375477 (1990), 2.374 (2010), 2.3728642 (2011), 2.3728639 (2014))

### 4 Master theorem (4.5 et 4.6)

$T(n) = aT(\frac{n}{b}) + f(n)$ . Let  $c = \log_b a$ .

1. If  $f(n) = O(n^{c'})$  with  $c' < c$ :  $T(n) = \Theta(n^c)$ .
  2. If  $f(n) = \Theta(n^c \log^k(n))$ :  $T(n) = \Theta(n^c \log^{k+1}(n))$ ;
  3. If  $f(n) = \Omega(n^{c'})$  with  $c' > c$  and if  $af(\frac{n}{b}) \leq \alpha f(n)$  for  $\alpha < 1$  with  $n$  large enough:  $T(n) = \Theta(f(n))$ .
- Not a trichotomy
  - Examples of uses
  - One still assumes  $n$  is a power of  $b$
  - Generalized by Akra–Bazzi (non-equal input size)
  - Proof! Essentially, Lemma 4.3 of textbook

### 5 Fast Polynomial Multiplication, Discrete Fourier Transform (30)

- Problem
- Naïve method,  $O(n^2)$
- Representation as sets of points, proof scheme
- $n$ th roots of 1
- Discrete Fourier transform (DFT) – connection with signal processing
- Fast Fourier transform: decomposition in even/odd components, DFT from these two decompositions, high-level algorithm scheme
- Complexity analysis with the master theorem
- Inverse DFT, Vandermonde matrix inverse, reuse of fast Fourier transform